

Pathfinder: XQuery—The Relational Way

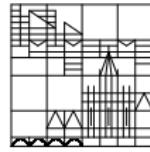
Jens Teubner¹ Peter Boncz² Torsten Grust¹
Maurice van Keulen³ Stefan Manegold² Jan Rittinger^{2,4}

¹Technische Universität München, Germany

²CWI Amsterdam, The Netherlands

³University of Twente, The Netherlands

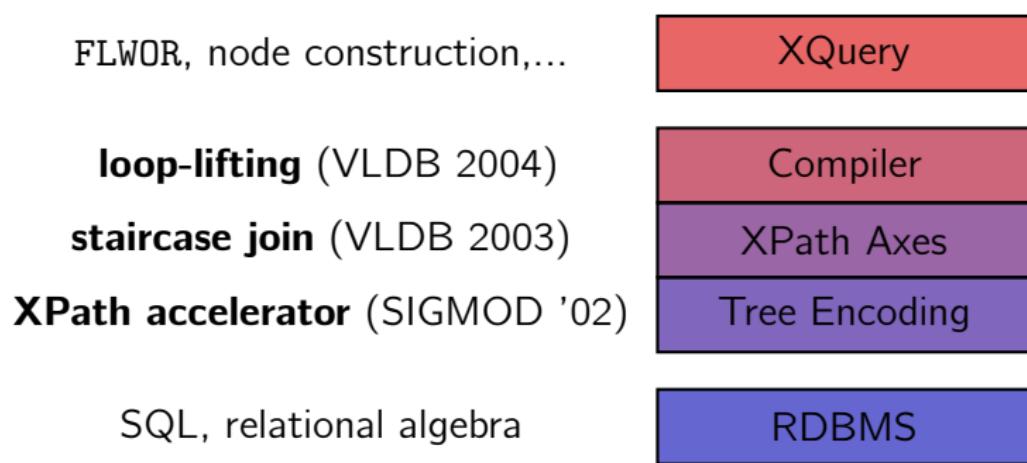
⁴University of Konstanz, Germany



Universität Konstanz

Pathfinder: Relational XQuery Processing

Pathfinder assembles relational XQuery processing techniques into a purely relational XQuery processing stack.



Source Language: XQuery Core

We provide full XQuery support.

literals	42, "foo", (), ...
arithmetics	$e_1 + e_2$, $e_1 - e_2$, ...
built-in functions	<code>fn:sum(e)</code> , <code>fn:count(e)</code> , <code>fn:doc(uri)</code> , ...
variable bindings	<code>let \$v := e₁ return e₂</code>
iteration	<code>for \$v at \$p in e₁ return e₂</code>
conditionals	<code>if p then e₁ else e₂</code>
sequence construction	e_1, e_2
function calls	$f(e_1, e_2, \dots, e_n)$
element construction	<code>element e₁ { e₂ }</code>
XPath steps	$e/\alpha::\nu$ (full axis feature)
:	:

- Expressions nest arbitrarily!

Target Language: Relational Algebra

A rather standard relational algebra suffices.

- π column projection, renaming
- σ row selection

Target Language: Relational Algebra

A rather standard relational algebra suffices.

π	column projection, renaming
σ	row selection
\bowtie	equi-join
\times	Cartesian product
$\dot{\cup}, \setminus$	disjoint union, difference
δ	duplicate elimination

Target Language: Relational Algebra

A rather standard relational algebra suffices.

π	column projection, renaming
σ	row selection
\bowtie	equi-join
\times	Cartesian product
$\dot{\cup}$, \	disjoint union, difference
δ	duplicate elimination
ϱ	row numbering

Target Language: Relational Algebra

A rather standard relational algebra suffices.

π	column projection, renaming
σ	row selection
\bowtie	equi-join
\times	Cartesian product
$\dot{\cup}, \setminus$	disjoint union, difference
δ	duplicate elimination
ϱ	row numbering
	staircase join*
ε, τ	element/text node construction*

*Syntactic sugar; expressible by remaining operators.

Target Language: Relational Algebra

A rather standard relational algebra suffices.

π	column projection, renaming
σ	row selection
\bowtie	equi-join
\times	Cartesian product
$\dot{\cup}, \setminus$	disjoint union, difference
δ	duplicate elimination
ϱ	row numbering
	staircase join*
ε, τ	element/text node construction*
\circledast	arithm./comparison operator *

*Syntactic sugar; expressible by remaining operators.

Target Language: Relational Algebra

A rather standard relational algebra suffices.

π	column projection, renaming
σ	row selection
\bowtie	equi-join
\times	Cartesian product
$\dot{\cup}, \setminus$	disjoint union, difference
δ	duplicate elimination
ϱ	row numbering
	staircase join*
ε, τ	element/text node construction*
\circledast	arithm./comparison operator *

- Operates on node (not tree!) level, 1NF relations.

*Syntactic sugar; expressible by remaining operators.

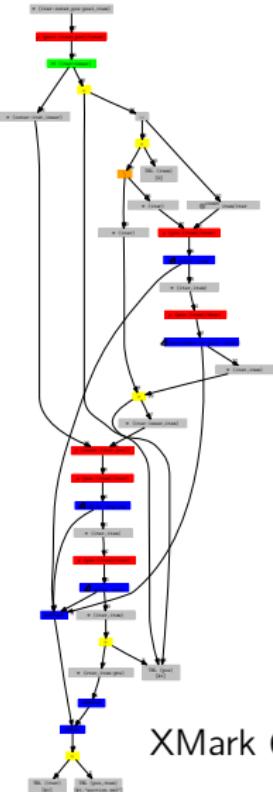
Target Language: Relational Algebra

A rather standard relational algebra suffices.

π	column projection, renaming
σ	row selection
\bowtie	equi-join
\times	Cartesian product
$\dot{\cup}, \setminus$	disjoint union, difference
δ	duplicate elimination
ϱ	row numbering
	staircase join*
ε, τ	element/text node construction*
\circledast	arithm./comparison operator *

- Operates on node (not tree!) level, 1NF relations.

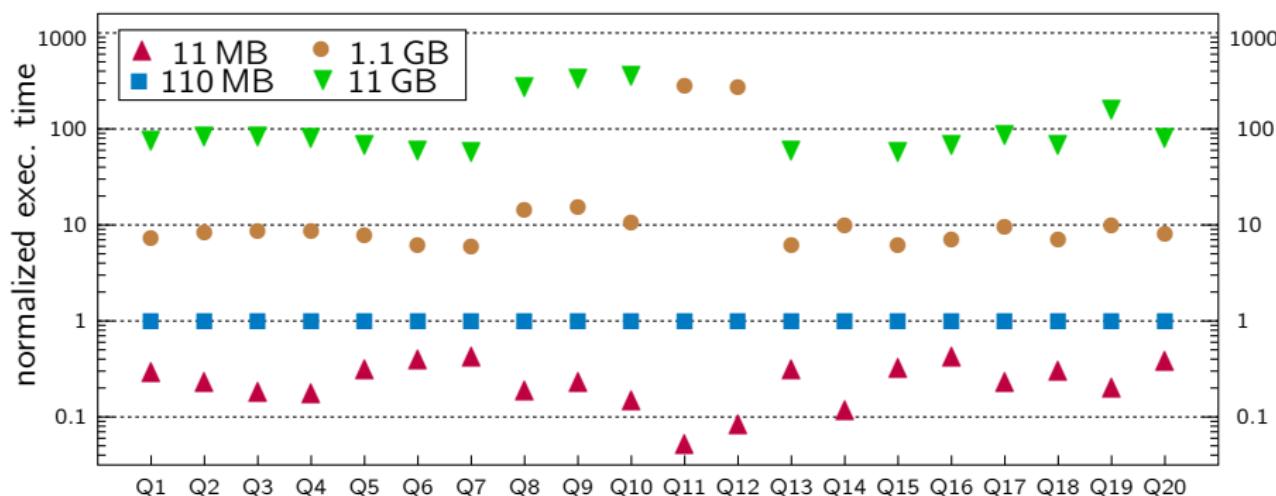
*Syntactic sugar; expressible by remaining operators.



Pathfinder Scalability

Version 0.8 of MonetDB/XQuery was released on May 30.

- Unsurpassed scalability, beyond 10 GB input document size.



Ongoing Work

The Pathfinder project is a joint effort of the Technische Universität München, CWI Amsterdam, and the University of Twente.

Optimizations:

- Algebraic join detection
- Order awareness (avoid ϱ)
- Use functional and multi-valued dependencies for algebraic optimization

Open Source Implementation:

<http://www.pathfinder-xquery.org/>

- Backed by main memory DBMS **MonetDB**.



See you in **Demo Group 7** (today, 16:00; tomorrow, 16:00)!